

DATABASE SYSTEM, METHOD FOR FORMING  
REPLICA OF DATABASE, AND COMPUTER-READABLE  
RECORDING MEDIUM THAT RECORDS DATABASE REPLICA FORMING PROGRAM

5 BACKGROUND OF THE INVENTION

This invention relates to forming of a replica of a database in a database system.

As described in chapter 23, pp. 624 and pp. 630 of "An Introduction to Database Systems, Volume I, Fifth Edition" by C. J. Date, the replica of database means a replica that should be identical with the master database in the data content and management method. A replica is prepared for avoiding a risk of working interruption due to work distribution of data processing or due to a failure of a computer that stores the data. It is required that the data content is logically identical between the master database and the replica. In order to be identical, it is required to apply the same data update processing as applied to the master database also to the replica to thereby maintain the same content. Heretofore, in order to maintain the same content, the same data update processing as applied to the master database is applied to the replica in the order of the occurrence of the update processing to thereby reflect the update on the replica.

A conventional replica forming process flow and the data content of update information to be used are shown in FIG. 6.

When replica forming process is started (401), for example, the update information used for updating the stored master database called as journal or update log is read in the order of occurrence successively (402). The read update information is edited so

5 as to be capable of reflecting on the replica side database (syntax or parameter format used to update the database), and then reflected on the replica side database (403). Whether there is any subsequent update information or not is determined (404), and if the result is YES, then the process is repeated.

10 When the update information reaches to the end of the stored update information, the sequence is brought to an end (405). The data content of the update information contains the update table name 501 that is the table name to be updated in the database, the type of update 502, and the record information 15 503 for indicating the updated record content. The table name is only an example, and the data element such as table, item, and record is the data content of the update information. The order of data of the update information data content of Fig. 6 is the same as the order of update that has occurred in the 20 master side database, and the update process of the replica occurs in the same order.

In the conventional art, the update process to be applied involves various tables (in other words, various data are mixed). Because update of the replica is applied in the same order as 25 applied actually in updating of the data even though usage and

necessity of the data are different respectively, it can occur that the necessary data is not reflected until all the stored update information is updated completely. For example, because update of another table is mixed during update of inventory table when the updated data of the inventory table is wanted in FIG. 6, it is not guaranteed that update of the inventory table has been completed until reflection of the update data is completed for all the update data. Therefore, it is a problem that it takes a long time to get the updated data that is required.

#### BRIEF SUMMARY OF THE INVENTION

In view of the above, it is the object of the present invention to get the updated data that is required preferentially out of various data correspondingly to the usage and necessity of the data.

In a database system having a master database and a replica thereof, the present invention involves reflection of the update information of the data on the replica not in the order of updating that has occurred in the master side but in the specified order of preference based on the purpose for using.

In detail, a memory unit for storing the order of preference of replica preparation process of a certain data type, an allocation processing unit for preferentially extracting the update data of the database correspondingly to the information, and a management unit for updating the replica successively in

the order according to the data extracted by means of the allocation processing means are provided.

#### BRIEF DESCRIPTION OF THE SEVERAL VIEWS OF THE DRAWING

5           FIG. 1 is a block diagram illustrating one example of the present invention.

FIG. 2 is a flow chart describing a data allocation unit shown in FIG. 1.

10           FIG. 3 is a diagram illustrating an example of preferential order specification in the example of the present invention.

FIG. 4 is a block diagram illustrating another example of the present invention.

15           FIG. 5 is a diagram illustrating an example of the database content and an example of preferential order information of the present invention.

FIG. 6 is a flow chart describing a conventional replica preparation process and an example of the update information data content.

20           FIG. 7 is a diagram illustrating another example of the present invention.

#### DETAILED DESCRIPTION OF THE PRERERRED EMBODIMENTS

Embodiments of the present invention will be described  
25 hereinafter with reference to the drawings. FIG. 1 is a

structural diagram of a database system of the present invention.

A master side database management unit 1 that is a system for managing a master database processes a data update request requested from a user, and updates a master DB 12 of a memory unit 5. The log information that has been got before and after updating is sent out to an update log file 13 when the data is updated for accommodating the failure of the master DB. A replica forming unit 3 reads an update log from the update log file 13 of the master side database, and updates the replica side database system. An update log reading unit 17 for reading an update log stores a necessary update data out of the read update logs in an update data cue 19.

A preferential order of the data, a replica of which is to be prepared, is specified by a use by use of a terminal unit key board. A preferential order acquiring unit 14 acquires the preferential order of the data, and holds it in a preferential order information memory unit 15 as the preferential order information. A data allocation unit 16 receives the preferential order information. The data allocation unit 16 reads out the information in the update data cue 19 successively in the order of updating, extracts the update information in the update data cue 19 in the order from the data type having the highest preference, and sends it out to a data update request unit 18. The data update request unit 18 sends out the received update request to a replica side database management unit 2.

The replica side database management unit 2 reflects the update request received from the data update request unit 18 on a replica DB 61.

The type of the data to which a preferential order is specified means, for example, the update information of the inventory table or merchandise master table in FIG. 6, or, for example, the update information of further subdivided category such as merchandise category A or merchandise category B in the inventory table. The preferential order may involve two-class classification in which the data type is classified into the highest order and others, or may involve multi-class classification in which the data type is classified into the first preferential order, the second preferential order, and so on. In the case of the two-class classification, not an order number but a flag is added simply to indicate the highest preference to be updated first. Otherwise, the data type to be updated preferentially may be stored.

FIG. 2 shows a process flow implemented when the replica forming unit 3 determines the preferential order and allocates the data. It is assumed that the preferential order is specified in the database table unit. When a data allocation process is started (101), the preferential order information of the table having the highest preference out of unprocessed tables is read based on the specified preference specifying information (102). Next, the update information is read from

the update data cue 19 (103), and the determination is made whether the table name having the update information is identical with the table name of the preferential order height that is being processed (104). If YES result is obtained, the  
5 update information is selected and sent to the data update request unit (105), and the replica side database management unit 2 reflects the update data on the replica.

As described herein above, the process in which the update information is read in the order of updating and the update  
10 information having the corresponding preferential order is selectively extracted is referred to as allocation. While the determination is made whether it is the end of the update information or not (106), and the update information is read and the process is repeated until end of the stored update  
15 information comes. When the end comes, whether there is the table having the next preferential order or not is determined (107), if there is the table having the next preferential order, then the same process as performed in the case is performed on the table.

20 As described hereinabove, the data allocation unit 16 selects the stored update information from that of the highest preferential order and transfers it to the data update request unit 18 for updating of the replica, the database is updated from the specified update information having the highest  
25 preferential order in the replica. As the result, the updated

information having higher preferential order request for using in the replica is rendered available earlier.

FIG. 3 shows an exemplary specification of the preferential order. In this example, three tables, namely the inventory data table 201, merchandise master table 202, and market analysis table 203, are extracted from the tables that are components of the master side database, and the update information is reflected on tables having the same names on the replica side. In the preference specification, the extraction side table 212 and reflection side table 213 are specified, and the preferential order 214 is specified finally. The result is stored in the preferential order information memory unit 15. In this exemplary specification, the updating of the inventory data 204 has the highest preference, and the updating of the market analysis 206 has the lowest preference.

FIG. 4 shows an exemplary system structure in which the replica forming unit is provided on the master side database and on the replica side database, both replica forming units are connected each other through a communication means, and the respective replica forming units specify the preference order. The master side replica forming unit is referred to as data extraction unit 31 and the replica side replica forming unit is referred to as data reflection unit 4. These replica forming units are somewhat different from the replica forming unit shown in FIG. 1 in the structure. The data extraction unit 31 and



the data reflection unit 4 are disposed separately, and the data is transferred between the data extraction unit 31 and the data reflection unit 4 by use of a communication means such as network, wireless communication network, communication cable, or

5 optical communication network. A data allocation unit 16 of the data extraction unit 31 sends the allocated data to a data transfer unit 28 in the same manner as shown in FIG. 1 and FIG. 2. The data transfer unit 28 sends the transferred update information to a data receiving unit 41 of the data reflection unit 4 by use of the communication means 7, and stores the update information in the update data cue 42.

An example in which the replica forming units disposed separately are used will be described with reference to FIG. 3 and FIG. 5 hereinafter. The exemplary preference specification shown in FIG. 3 is used as the preference specification for the data extraction unit 31. The inventory data shown in FIG. 5 is an exemplary inventory data table stored in the database 12. The preference specification a and the preference specification b are different examples of the preference specification, and stored in a preferential order information memory unit 52. The necessity preference order of the data, namely the inventory data> the merchandise master> the market analysis, is set in the replica side system, and the preference order acquiring unit 14 of the data extraction unit 31 acquires the preference order.

The preference specification in the data reflection unit 4 is shown in FIG. 5. The preference specification a is an example in which the record that preferentially reflects the update information of the inventory data based on the

5 merchandise category 301 is selected. The preferential key condition 312 and the preferential order 313 that are to be referred now is specified out of the inventory data on the replica side to thereby not only preferentially specify the inventory data but also further subdivide the preferential order specification. Herein, the data is updated to the newest  
10 information for respective merchandise category data of the inventory data. According to the preferential order, because the merchandise category reflects the inventory of the merchandise A preferentially among the inventory data, the  
15 preferential reflection is used effectively when the inventory of the merchandise category A is to be checked during business negotiation.

The preference specification b shows an example in which a necessary item is reflected preferentially. When the  
20 database is used for other purpose such as data analysis, only the item that is necessary for analysis is reflected preferentially and the residual items are reflected later. As the result, the analysis can be started earlier and the earlier starting is effective in the case of the time-consuming analysis  
25 such as multidimensional analysis. In the preference

specification b, the preferential column 322 of the table of the inventory data is specified and the preferential order 323 is specified respectively to thereby reflect the merchandise category, merchandise code, and inventory preferentially.

5           Examples, in which the update data is extracted preferentially for the category including to the table name on the master side, and further subdivided preferential key condition is specified and the update data is extracted preferentially by specifying the preferential column of the  
10 table, are shown hereinbefore. According to the above, the database can be updated based on the subdivided preferential order and also the following effect is obtained. In detail, it is assumed that a plurality of data reflection units 4 and replicas are provided and only the inventory data is transferred  
15 from the data transfer unit 28 of the data extraction unit 31. A certain data reflection unit allocates the merchandise category A as the most preferential update data, and another data reflection unit allocates the merchandise category C as the most preferential update data. The replica is  
20 preferentially updated so as to satisfy different requests requested from users located on remote areas distant far from the master DB by employing the structure as described hereinabove.

Furthermore, another example is described herein under.

25 The preferential order acquiring unit 14, the preferential order

information memory unit 15, and the data allocation unit 16 of the data extraction unit 31 shown in FIG. 4 are removed or not used, and all update data are transferred to the data reflection unit 4 on the replica side through the communication means 7.

- 5 The allocation preferential order of the update data is determined only by means of the data reflection unit 4. Thereby, the master DB side load is reduced.

Another example, in which a plurality of data reflection units 4 and replicas are provided as shown in FIG. 4 and on the other hand the data extraction unit 31 extracts a plurality of types of update data in parallel, will be described herein under. The plurality of types of update data are transferred to different data reflection units 4 respectively by means of the data transfer unit 28. Thereby, the replica can be updated so as to satisfy the request of users located on remote areas.

FIG. 7 shows a diagram illustrating another example of the present invention. In this example, the replica side database management unit 2 has acquired the database reference history information, and the update information of the table of the database that has been referred most frequently is reflected preferentially on the replica. Therefore, the preferential order is specified not based on the user specification but based on the result of statistical processing of actual database reference history. Differently from the example shown in FIG. 4, what table of the replica has been

referred in the replica side database management unit 2 is transferred to the replica reference history acquiring unit 55, and the replica reference history acquiring unit 55 sets it to the preferential order information memory unit 52 as the

5 preferential order information. The data allocation unit 43 develops the update information on the data update request unit 44 by use of the information supplied from the preferential order information memory unit 52. Thereby, it is possible that the system judges the preferential order on which the update  
10 information is reflected according to the reference frequency of the replica side database used by users and the data of high reference frequency is reflected preferentially as the newest data.

According to the present invention, the update order is  
15 changed so that the more necessary data is updated earlier and less necessary data is updated later by means of the allocation unit. However, after all, the update data of all the data types is reflected on the replica in the same order as the update order of the master database. Some data type is updated earlier  
20 according the specified preferential order.

Examples in which the database holds the data in the form of table are described hereinbefore, but the present invention is by no means limited to the specified database type. The above-mentioned process is performed by means of a program in  
25 the examples. The program is recorded in a computer-readable

According to the present invention, it is possible to reflect the update information with selection of the data type, the replica of which is to be prepared preferentially

NAME	ADDRESS	CITY	STATE	ZIP
Mr. J. H. Smith	123 Main St.	Springfield	Ill.	62761
Mr. W. R. Jones	456 Oak Ave.	Chicago	Ill.	60601
Mr. T. L. Brown	789 Elm St.	Peoria	Ill.	61601
Mr. S. K. White	101 Maple Dr.	Rockford	Ill.	61101
Mr. D. M. Green	202 Pine St.	Decatur	Ill.	62521
Mr. C. N. Black	303 Cedar Ave.	Normal	Ill.	62451
Mr. E. O. Gray	404 Birch St.	Urbana	Ill.	61501
Mr. F. P. Hall	505 Walnut Dr.	Champaign	Ill.	61821
Mr. G. Q. King	606 Spruce Ave.	Carbondale	Ill.	62901
Mr. H. R. Lee	707 Fir St.	Macomb	Ill.	61451
Mr. I. S. Walker	808 Ash Dr.	Edwardsville	Ill.	62021
Mr. J. T. Young	909 Hickory Ave.	St. Louis	Mo.	63101
Mr. K. U. Allen	1010 Sycamore St.	St. Charles	Mo.	63071
Mr. L. V. Wright	1111 Chestnut Dr.	St. Joseph	Mo.	64501
Mr. M. W. Scott	1212 Locust Ave.	St. Louis	Mo.	63101
Mr. N. X. Adams	1313 Walnut St.	St. Louis	Mo.	63101
Mr. O. Y. Baker	1414 Elm Ave.	St. Louis	Mo.	63101
Mr. P. Z. Carter	1515 Oak St.	St. Louis	Mo.	63101
Mr. Q. A. Evans	1616 Maple Dr.	St. Louis	Mo.	63101
Mr. R. B. Foster	1717 Pine St.	St. Louis	Mo.	63101
Mr. S. C. Gibson	1818 Cedar Ave.	St. Louis	Mo.	63101
Mr. T. D. Hall	1919 Birch St.	St. Louis	Mo.	63101
Mr. U. E. King	2020 Walnut Dr.	St. Louis	Mo.	63101
Mr. V. F. Lee	2121 Spruce Ave.	St. Louis	Mo.	63101
Mr. W. G. Miller	2222 Fir St.	St. Louis	Mo.	63101
Mr. X. H. Wilson	2323 Ash Dr.	St. Louis	Mo.	63101
Mr. Y. I. Moore	2424 Hickory Ave.	St. Louis	Mo.	63101
Mr. Z. J. Taylor	2525 Sycamore St.	St. Louis	Mo.	63101
Mr. A. K. Anderson	2626 Chestnut Dr.	St. Louis	Mo.	63101
Mr. B. L. Jackson	2727 Locust Ave.	St. Louis	Mo.	63101
Mr. C. M. Roberts	2828 Walnut St.	St. Louis	Mo.	63101
Mr. D. N. Turner	2929 Elm Ave.	St. Louis	Mo.	63101
Mr. E. O. Phillips	3030 Oak St.	St. Louis	Mo.	63101
Mr. F. P. Campbell	3131 Maple Dr.	St. Louis	Mo.	63101
Mr. G. Q. Parker	3232 Pine St.	St. Louis	Mo.	63101
Mr. H. R. Evans	3333 Cedar Ave.	St. Louis	Mo.	63101
Mr. I. S. King	3434 Birch St.	St. Louis	Mo.	63101
Mr. J. T. Lee	3535 Walnut Dr.	St. Louis	Mo.	63101
Mr. K. U. Miller	3636 Spruce Ave.	St. Louis	Mo.	63101
Mr. L. V. Wilson	3737 Fir St.	St. Louis	Mo.	63101
Mr. M. W. Moore	3838 Ash Dr.	St. Louis	Mo.	63101
Mr. N. X. Taylor	3939 Hickory Ave.	St. Louis	Mo.	63101
Mr. O. Y. Anderson	4040 Sycamore St.	St. Louis	Mo.	63101
Mr. P. Z. Jackson	4141 Chestnut Dr.	St. Louis	Mo.	63101
Mr. Q. A. Roberts	4242 Locust Ave.	St. Louis	Mo.	63101
Mr. R. B. Turner	4343 Walnut St.	St. Louis	Mo.	63101
Mr. S. C. Phillips	4444 Elm Ave.	St. Louis	Mo.	63101
Mr. T. D. Campbell	4545 Oak St.	St. Louis	Mo.	63101
Mr. U. E. Parker	4646 Maple Dr.	St. Louis	Mo.	63101
Mr. V. F. Evans	4747 Pine St.	St. Louis	Mo.	63101
Mr. W. G. King	4848 Cedar Ave.	St. Louis	Mo.	63101
Mr. X. H. Lee	4949 Birch St.	St. Louis	Mo.	63101
Mr. Y. I. Miller	5050 Walnut Dr.	St. Louis	Mo.	63101
Mr. Z. J. Wilson	5151 Spruce Ave.	St. Louis	Mo.	63101
Mr. A. K. Moore	5252 Fir St.	St. Louis	Mo.	63101
Mr. B. L. Taylor	5353 Ash Dr.	St. Louis	Mo.	63101
Mr. C. M. Anderson	5454 Hickory Ave.	St. Louis	Mo.	63101
Mr. D. N. Jackson	5555 Sycamore St.	St. Louis	Mo.	63101
Mr. E. O. Roberts	5656 Chestnut Dr.	St. Louis	Mo.	63101
Mr. F. P. Turner	5757 Locust Ave.	St. Louis	Mo.	63101
Mr. G. Q. Phillips	5858 Walnut St.	St. Louis	Mo.	63101
Mr. H. R. Campbell	5959 Elm Ave.	St. Louis	Mo.	63101
Mr. I. S. Parker	6060 Oak St.	St. Louis	Mo.	63101
Mr. J. T. Evans	6161 Maple Dr.	St. Louis	Mo.	63101
Mr. K. U. King	6262 Pine St.	St. Louis	Mo.	63101
Mr. L. V. Lee	6363 Cedar Ave.	St. Louis	Mo.	63101
Mr. M. W. Miller	6464 Birch St.	St. Louis	Mo.	63101
Mr. N. X. Wilson	6565 Walnut Dr.	St. Louis	Mo.	63101
Mr. O. Y. Moore	6666 Spruce Ave			